# First-Order Logic Part One 

- A propositional variable is a variable that is either true or false.
- The propositional connectives are as follows:
- Negation: $\neg p$
- Conjunction: $p \wedge q$
- Disjunction: $p \vee q$
- Implication: $p \rightarrow q$
- Biconditional: $p \leftrightarrow q$
- True: T
- False: $\perp$


## What is First-Order Logic?

- First-order logic is a logical system for reasoning about properties of objects.
- Augments the logical connectives from propositional logic with
- predicates that describe properties of objects,
- functions that map objects to one another, and
- quantifiers that allow us to reason about multiple objects.


## Some Examples

$\operatorname{Likes}($ You, Eggs) $\wedge \operatorname{Likes(You,~Tomato)~} \rightarrow \operatorname{Likes(You,~Shakshuka)~}$

$\operatorname{Likes}($ You, Eggs) $\wedge \operatorname{Likes(You,~Tomato)~} \rightarrow$ Likes(You, Shakshuka)
Learns(You, History) v ForeverRepeats(You, History)
In(MyHeart, Havana) ^ TookBackTo(Him, Me, EastAtlanta)

Likes(You, Eggs) ^Likes(You, Tomato) $\rightarrow$ Likes (You, Shakshuka)
Learns(You, History) v ForeverRepeats(You, History)
In(MyHeart, Havana) ^ TookBackTo(Him, Me, EastAtlanta)

## Likes(You, Eggs) ^Likes(You, Tomato) $\rightarrow$ Likes(You, Shakshuka)

 Learns(You, History) v ForeverRepeats(You, History)In(MyHeart, Havana) ^ TookBackTo(Him, Me, EastAtlanta)

These blue terms are called constant symbols. Unlike propositional variables, they refer to objects, not propositions.
$\operatorname{Likes}($ You, Eggs $) \wedge \operatorname{Likes}($ You, Tomato $) \rightarrow \operatorname{Likes}($ You, Shakshuka)
Learns(You, History) v ForeverRepeats(You, History)
In(MyHeart, Havana) ^ TookBackTo(Him, Me, EastAtlanta)
$\operatorname{Likes}($ You, Eggs $) \wedge \operatorname{Likes}($ You, Tomato $) \rightarrow \operatorname{Likes}($ You, Shakshuka)
Learns(You, History) v ForeverRepeats(You, History)
In(MyHeart, Havana) ^ TookBackTo(Him, Me, EastAtlanta)

The red things that look like function calls are called predicates. Predicates take objects as arguments and evaluate to true or false.

Likes(You, Eggs) ^Likes(You, Tomato) $\rightarrow$ Likes(You, Shakshuka)
Learns(You, History) v ForeverRepeats(You, History)
In(MyHeart, Havana) ^ TookBackTo(Him, Me, EastAtlanta)

# Likes(You, Eggs) ^ Likes(You, Tomato) $\rightarrow$ Likes(You, Shakshuka) 

Learns(You, History) v ForeverRepeats(You, History)
In(MyHeart, Havana) ^ TookBackTo(Him, Me, EastAtlanta)

What remains are traditional propositional connectives. Because each predicate evaluates to true or false, we can connect the truth values of predicates using normal propositional connectives.

## Reasoning about Objects

- To reason about objects, first-order logic uses predicates.
- Examples:

Cute(Quokka)
ArgueIncessantly(Democrats, Republicans)

- Applying a predicate to arguments produces a proposition, which is either true or false.
- Typically, when you're working in FOL, you'll have a list of predicates, what they stand for, and how many arguments they take. It'll be given separately than the formulas you write.


## First-Order Sentences

- Sentences in first-order logic can be constructed from predicates applied to objects:

$$
\begin{gathered}
\text { Cute }(a) \rightarrow \operatorname{Dikdik}(a) \vee \operatorname{Kitty}(a) \text { v Puppy }(a) \\
\text { Succeeds }(Y o u)
\end{gathered} \leftrightarrow \operatorname{Practices(You)} \text { (Y) }
$$

$$
x<8 \rightarrow x<137
$$

The less -than sign is just another predicate. Binary predicates are sometimes written in infix notation this way.

Numbers are not "built
$\mathrm{in}^{\text {" }}$ to first-order logic. They're constant symbols just like "You" and "a" above.

## Equality

- First-order logic is equipped with a special predicate $=$ that says whether two objects are equal to one another.
- Equality is a part of first-order logic, just as $\rightarrow$ and $\neg$ are.
- Examples:

$$
\begin{gathered}
\text { TomMarvoloRiddle }=\text { LordVoldemort } \\
\text { MorningStar }=\text { EveningStar }
\end{gathered}
$$

- Equality can only be applied to objects; to state that two propositions are equal, use $\leftrightarrow$.


## Let's see some more examples.

FavoriteMovieOf(You) $\neq$ FavoriteMovieOf(Date) $\wedge$ StarOf(FavoriteMovieOf(You)) = StarOf(FavoriteMovieOf(Date))

FavoriteMovieOf(You) $\neq$ FavoriteMovieOf(Date) $\wedge$ StarOf(FavoriteMovieOf(You)) $=$ StarOf(FavoriteMovieOf(Date))

# FavoriteMovieOf(You) $\neq$ FavoriteMovieOf(Date) ^ StarOf(FavoriteMovieOf(You)) $=$ StarOf(FavoriteMovieOf(Date)) 

FavoriteMovieOf(You) $\neq$ FavoriteMovieOf(Date) $\wedge$ StarOf(FavoriteMovieOf(You)) $=$ StarOf(FavoriteMovieOf(Date))

# FavoriteMovieOf(You) $\neq$ FavoriteMovieOf(Date) ^ StarOf(FavoriteMovieOf(You)) = StarOf(FavoriteMovieOf(Date)) 

These purple terms are functions. Functions take objects as input and produce objects as output.

FavoriteMovieOf(You) $\neq$ FavoriteMovieOf(Date) $\wedge$ StarOf(FavoriteMovieOf(You)) $=$ StarOf(FavoriteMovieOf(Date))

FavoriteMovieOf(You) $\neq$ FavoriteMovieOf(Date) $\wedge$ StarOf(FavoriteMovieOf(You)) $=$ StarOf(FavoriteMovieOf(Date))

FavoriteMovieOf(You) $\neq$ FavoriteMovieOf(Date) $\wedge$ StarOf(FavoriteMovieOf(You)) $=$ StarOf(FavoriteMovieOf(Date))

## Functions

- First-order logic allows functions that return objects associated with other objects.
- Examples:

$$
\begin{gathered}
\text { ColorOf(Money) } \\
\text { MedianOf( } x, y, z) \\
x+y
\end{gathered}
$$

- As with predicates, functions can take in any number of arguments, but always return a single value.
- Functions evaluate to objects, not propositions.


## Objects and Predicates

- When working in first-order logic, be careful to keep objects (actual things) and propositions (true or false) separate.
- You cannot apply connectives to objects: © Venus $\rightarrow$ TheSun
- You cannot apply functions to propositions:
© StarOf(IsRed(Sun) ^ IsGreen(Mars)) ©
- Ever get confused? Just ask!


## Some muggle is intelligent.

## Some muggle is intelligent.

$\exists m$. (Muggle(m) ^ Intelligent(m))

## Some muggle is intelligent.

## $\exists m .(M u g g l e(m) \wedge$ Intelligent (m))

$\exists$ is the existential quantifier and says "for some choice of $m$, the following is true."

## The Existential Quantifier

- A statement of the form

$$
\exists x . \text { some-formula }
$$

is true if, for some choice of $x$, the statement some-formula is true when that $x$ is plugged into it.

- Examples:
$\exists x .(E v e n(x) \wedge \operatorname{Prime}(x))$
$\exists x$. (TallerThan(x, me) $\wedge$ LighterThan(x, me))
$(\exists w . \operatorname{Will}(w)) \rightarrow(\exists x . \operatorname{Way}(x))$


## The Existential Quantifier


$\exists x . \operatorname{Smiling}(x)$

## The Existential Quantifier



## The Existential Quantifier



## The Existential Quantifier



## The Existential Quantifier



## The Existential Quantifier



## The Existential Quantifier



## The Existential Quantifier



## The Existential Quantifier



## The Existential Quantifier


$\exists x . \operatorname{Smiling}(x)$

## The Existential Quantifier



## The Existential Quantifier



## The Existential Quantifier



## The Existential Quantifier



## The Existential Quantifier



## The Existential Quantifier



## The Existential Quantifier



## The Existential Quantifier



## The Existential Quantifier


$(\exists x . S m i l i n g(x)) \rightarrow(\exists y$. WearingHat(y))

## The Existential Quantifier


$(\exists x . \operatorname{Smiling}(x)) \rightarrow(\exists y$. WearingHat(y))

## The Existential Quantifier



Is this part of the statement true or false?
$(\exists x . \operatorname{Smiling}(x)) \rightarrow(\exists y$. WearingHat(y))

## The Existential Quantifier



Is this part of the statement true or false?
$(\exists x . \operatorname{Smiling}(x)) \rightarrow(\exists y$. WearingHat $(y))$

## The Existential Quantifier



Is this part of the statement true or false?
$(\exists x . S m i l i n g(x)) \rightarrow(\exists y$. WearingHat(y))

## The Existential Quantifier



Is this part of the statement true or false?
$(\exists x . \operatorname{Smiling}(x)) \rightarrow(\exists y$. WearingHat(y))

## The Existential Quantifier


$(\exists x . \operatorname{Smiling}(x)) \rightarrow(\exists y$. WearingHat $(y))$

## The Existential Quantifier


$(\exists x$. Smiling $(x)) \rightarrow(\exists y$. WearingHat $(y))$

## Fun with Edge Cases

$\exists x . \operatorname{Smiling}(x)$

## Fun with Edge Cases

Existentially-quantified statements are false in an empty world, since nothing exists, period!
$\exists x$. Smiling $(x)$
"For any natural number $n$, $n$ is even if and only if $n^{2}$ is even"
$\forall n .\left(n \in \mathbb{N} \rightarrow\left(\operatorname{Even}(n) \leftrightarrow \operatorname{Even}\left(n^{2}\right)\right)\right)$
"For any natural number $n$, $n$ is even if and only if $n^{2}$ is even"

## $\forall n .\left(n \in \mathbb{N} \rightarrow\left(\operatorname{Even}(n) \leftrightarrow \operatorname{Even}\left(n^{2}\right)\right)\right)$

$\forall$ is the universal quantifier and says "for any choice of $n$, the following is true."

## The Universal Quantifier

- A statement of the form


## $\forall x$. some-formula

is true if, for every choice of $x$, the statement some-formula is true when $x$ is plugged into it.

- Examples:
$\forall$. $(\operatorname{Puppy}(p) \rightarrow$ Cute $(p))$
$\forall a$. (EatsPlants(a) v EatsAnimals(a))
Tallest(SultanKösen) $\rightarrow$
$\forall x$. (SultanKösen $\neq x \rightarrow$ ShorterThan( $x$, SultanKösen))


## The Universal Quantifier


$\forall x . \operatorname{Smiling}(x)$

## The Universal Quantifier


$\forall x . \operatorname{Smiling}(x)$

## The Universal Quantifier



## The Universal Quantifier



## The Universal Quantifier



## The Universal Quantifier



## The Universal Quantifier



## The Universal Quantifier



Since Smiling ( $x$ ) is true for every choice of $x$, this statement evaluates to true.
$\forall x . \operatorname{Smiling}(x)$

## The Universal Quantifier



Since Smiling ( $x$ ) is true for every choice of $x$, this statement evaluates to true.
$\forall x . \operatorname{Smiling}(x)$

## The Universal Quantifier


$\forall x . \operatorname{Smiling}(x)$

## The Universal Quantifier


$\forall x . \operatorname{Smiling}(x)$

## The Universal Quantifier



## The Universal Quantifier



## The Universal Quantifier



Since Smiling(x) is false for this choice $x$, this statement evaluates to false.

## The Universal Quantifier



Since Smiling(x) is false for this choice $x$, this statement evaluates to false.

## The Universal Quantifier


$(\forall x . \operatorname{Smiling}(x)) \rightarrow(\forall y$. WearingHat $(y))$

## The Universal Quantifier


$(\forall x . S m i l i n g(x)) \rightarrow(\forall y$. WearingHat $(y))$

## The Universal Quantifier



Is this part of the statement true or false?
$(\forall x$. Smiling $(x)) \rightarrow(\forall y$. WearingHat $(y))$

## The Universal Quantifier



Is this part of the statement true or false?
$(\forall x$. Smiling $(x)) \rightarrow(\forall y$. WearingHat $(y))$

## The Universal Quantifier



Is this part of the statement true or false?
$(\forall x$. Smiling $(x)) \rightarrow(\forall y$. WearingHat $(y))$

## The Universal Quantifier



Is this part of the statement true or false?
$(\forall x$.Smiling $(x)) \rightarrow(\forall y$. WearingHat $(y))$

## The Universal Quantifier



Is this overall statement true or false in this scenario?
$(\forall x$. Smiling $(x)) \rightarrow(\forall y$. WearingHat $(y))$

## The Universal Quantifier



Is this overall statement true or false in this scenario?

## $(\forall x . \operatorname{Smiling}(x)) \rightarrow(\forall y$. WearingHat $(y))$

